Synthesis of asynchronous reactive programs from temporal specifications

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Hotel reservation



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Designing asynchronous systems is hard



Partial/Outdated information



Unpredictable delays

Can we automatically generate an asynchronous program from its specification?

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Asynchronous synthesis

Can we automatically generate an asynchronous program from its specification written in LTL?

Asynchronous LTL synthesis

2EXPTIME-complete [Pnueli, Rosner; ICALP 1989] But no tools!

Towards building tools/techniques for asynchronous LTL synthesis

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Simpler theoretical analysis

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First tool for asynchronous synthesis for all of LTL Existing tools restricted to specific LTL formula [Klein, Piterman, Pnueli; VMCAI 2012] Direct translation to SAT causes blowup [Schewe, Finkbeiner; STTT 2013]

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Identify LTL subclasses with lower complexity

[Pnueli, Rosner; ICALP 1989]

- Reactive system
- Receives input sequence I, observes a sub-sequence X only
- Generates output sequence **O** based on **X** only
- Outputs generated with (non-fixed) delay

[Pnueli, Rosner; ICALP 1989]

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Asynchronous LTL synthesis [Pnueli, Rosner; ICALP 1989]

Given LTL specification **S**, Generate asynch. system **f** s.t. (I,O) satisfies **S** for all I and **X**

Challenge I: Partial information

Asych. system satisfies specification despite not reading all inputs

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Asych. system satisfies specification despite not reading all inputs

Example: Always (i \rightarrow Eventually o)



Input (I)	Т	Т	F	Т	F	F	Т	•••
Output (O)	-	-	-	F	F	F	F	•••

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Asych. system satisfies specification despite not reading all inputs

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Output (O)	-	-	-	Т	Т	Т	Т	•••







Between observing input and generating output



(I,O) satisfies **S** iff all expansions and compressions of (I,O) satisfy **S**

Previous work [Pnueli, Rosner; ICALP 1989]

A trace satisfies **TS** iff all expansion and compression of the trace satisfy **S**



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Our reduction: Pnueli-Rosner (PR) closure

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Implementation and Experiments

- Prototype tool BAS (Bounded Asynchronous Synthesis)
 - BDD-based symbolic construction of **PR(S)** from **S**
 - Bounded synch. synthesis tool
 - **BoSy** [Faymonville, Finkbeiner, Tentrup; CAV 2017]
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 - Number of propositions in range of 2-16
 - Length of LTL-specification ranges from 5-80

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- Evaluate performance on ~15 formulae
 - Number of propositions in range of 2-16
 - Length of LTL-specification ranges from 5-80
 - Almost all formulae solved in less than 2 second
 - $PR(\varphi)$ construction takes ~2% of total time

Asynchronous LTL synthesis, in a nutshell

- Efficient reduction of asynch. LTL synthesis to synch. synthesis
 - Modified spec. PR closure is exponentially smaller than previous reduction
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- Low-complexity subclasses of GR(1) spec.
 - Reduction to single-alternation QBF satisfiability
 - Much lower complexity than the general 2EXPTIME